



# An Approach for Encryption and Decryption Using Matrix Theory

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## Abstract

In this paper, a method for encryption and decryption have been discussed using matrix theory and we have illustrated an example for this method.

**Key words:** Encryption, Decryption, Matrix, Key Matrix.

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## 1. Introduction

Cryptography is the research and practise of safe communication strategies in the presence of adversaries identified by third parties. It deals with the development and analysis of protocols that prevent malicious third parties from accessing information that is shared between two entities, thus pursuing the different aspects of information security. The encryption process is comprised of a Algorithm, with a key [4]. The ordinary readable text is known as plain text. The transformed message is known as cipher text. The key is the value independent of the plaintext.

The conversion process of plain text to cipher text is called encryption algorithm and cipher text to plain text is called decryption algorithm. In [1], they have developed a technique for encryption and decryption using matrix theory.

## 2. Mathematical concepts

**Theorem 2.1** A text message of strings of some length size  $L$  can be converted into a matrix called a message matrix  $R$  of size  $n > m$  and  $n$  is the least such that  $m \times n \geq L$  depending upon the length of the message with the help of suitably chosen numerical and zeros [3].

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### 3. Algorithms

#### Encryption Algorithm:

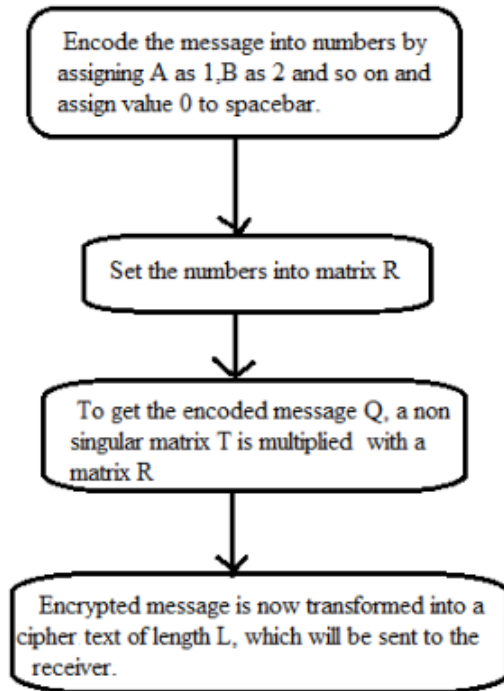


Figure 1: Procedure for Encoding

#### Decryption Algorithm :

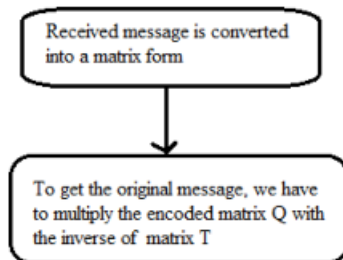


Figure 2: Procedure for Decoding

### 4. Results and Discussions

**Illustration 4.1** The message which we are going to send to the receiver is *MEET ME AT BUTTERFLY GALAXY TO GET HONEY*

Now let us convert the message into numbers ,

**13 5 5 20 0 13 5 0 1 20 0 2 21 20 20 5 18 6 12 25 0 7**  
**1 12 1 24 25 0 20 15 0 7 5 20 0 8 15 14 5 25**

Setting these numbers into a matrix form R as below,

$$R = \begin{bmatrix} 13 & 5 & 5 & 20 & 0 \\ 13 & 5 & 0 & 1 & 20 \\ 0 & 2 & 21 & 20 & 20 \\ 5 & 18 & 6 & 12 & 25 \\ 0 & 7 & 1 & 12 & 1 \\ 24 & 25 & 0 & 20 & 15 \\ 0 & 7 & 5 & 20 & 0 \\ 8 & 15 & 14 & 5 & 25 \end{bmatrix}$$

Now let us assume a non singular matrix T as,  $T = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$

as an encryption key then  $T^{-1}$  is  $\begin{bmatrix} 1 & -1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$

We now multiplied matrix R with a non singular matrix T to get the encoded matrix Q .

$$Q = RT = \begin{bmatrix} 13 & 5 & 5 & 20 & 0 \\ 13 & 5 & 0 & 1 & 20 \\ 0 & 2 & 21 & 20 & 20 \\ 5 & 18 & 6 & 12 & 25 \\ 0 & 7 & 1 & 12 & 1 \\ 24 & 25 & 0 & 20 & 15 \\ 0 & 7 & 5 & 20 & 0 \\ 8 & 15 & 14 & 5 & 25 \end{bmatrix} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$Q = \begin{bmatrix} 13 & 18 & 23 & 43 & 43 \\ 13 & 18 & 18 & 19 & 39 \\ 0 & 2 & 23 & 43 & 63 \\ 5 & 23 & 29 & 41 & 66 \\ 0 & 7 & 8 & 20 & 21 \\ 24 & 49 & 49 & 69 & 84 \\ 0 & 7 & 12 & 32 & 32 \\ 8 & 23 & 37 & 42 & 67 \end{bmatrix}.$$

The encoded message to be sent is

***13 18 23 43 43 13 18 18 19 39 0 2 23 43 63 5 23 29 41 66  
 0 7 8 20 21 24 49 49 69 84 0 7 12 32 32 8 23 37 42 67.***

To get the original message receiver should multiply by  $T^{-1}$

$$R = RTT^{-1} = \begin{bmatrix} 13 & 18 & 23 & 43 & 43 \\ 13 & 18 & 18 & 19 & 39 \\ 0 & 2 & 23 & 43 & 63 \\ 5 & 23 & 29 & 41 & 66 \\ 0 & 7 & 8 & 20 & 21 \\ 24 & 49 & 49 & 69 & 84 \\ 0 & 7 & 12 & 32 & 32 \\ 8 & 23 & 37 & 42 & 67 \end{bmatrix} \begin{bmatrix} 1 & -1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Therefore, the decoded message is,

**13 5 5 20 0 13 5 0 1 20 0 2 21 20 20 5 18 6 12 25 0 7  
 1 12 1 24 25 0 20 15 0 7 5 20 0 8 15 14 5 25**

Hence, we received the original plaintext by changing the numbers into alphabets. We get the original message as **MEET ME AT BUTTERFLY GALAXY TO GET HONEY**

**NOTE:** We have used Matlab for matrix multiplication.

## 5. Congruence Modulo Method

**Definition 5.1** Let  $r$  be a positive integer, we say that  $u$  is congruent to  $v \pmod{r}$  if  $r|(u-v)$  where  $u$  and  $v$  are integers i.e.,  $u = v + sr$  and  $s \in \mathbb{Z}$ , we write  $u \equiv v \pmod{r}$  is called congruence relation, the number  $r$  is the modulus of congruence [2],[5].

**Definition 5.2** Inverse of an integer  $h$  to modulo  $r$  is  $h^{-1}$  such that  $[h.h]^{-1} \equiv 1 \pmod{r}$ , where  $h^{-1}$  is called inverse of  $h$ .

**Illustration 5.3** First we are going to assign numbers from 1 to 26 to the 26 alphabets starting from A to Z and 0 to spacebar. Since we are going to use congruence method so let us take matrix modulo 27. Consider the message that is plain text as **MEET ME AT BUTTERFLY GALAXY TO GET HONEY**

Alphabet	A	B	C	D	E	F	G	H	I
Number	1	2	3	4	5	6	7	8	9
	-26	-25	-24	-23	-22	-21	-20	-19	-18
Alphabet	J	K	L	M	N	O	P	Q	R
Number	10	11	12	13	14	15	16	17	18
	-17	-16	-15	-14	-13	-12	-11	-10	-9
Alphabet	S	T	U	V	W	X	Y	Z	spacebar
Number	19	20	21	22	23	24	25	26	0
	-8	-7	-6	-5	-4	-3	-2	-1	0

Now let us assign the numbers to the above words by using above table , and we are going to arrange it in  $5 \times 1$  matrix.

$$\begin{aligned}
 \text{MEET} &= \begin{bmatrix} 13 \\ 5 \\ 5 \\ 20 \\ 0 \end{bmatrix} ; \text{ME AT} = \begin{bmatrix} 13 \\ 5 \\ 0 \\ 1 \\ 20 \end{bmatrix} ; \text{BUTT} = \begin{bmatrix} 0 \\ 2 \\ 21 \\ 20 \\ 20 \end{bmatrix} ; \text{ERFLY} = \begin{bmatrix} 5 \\ 18 \\ 6 \\ 12 \\ 25 \end{bmatrix} ; \\
 \text{GALA} &= \begin{bmatrix} 0 \\ 7 \\ 1 \\ 12 \\ 1 \end{bmatrix} ; \text{XY TO} = \begin{bmatrix} 24 \\ 25 \\ 0 \\ 20 \\ 15 \end{bmatrix} ; \text{GET} = \begin{bmatrix} 0 \\ 7 \\ 5 \\ 20 \\ 0 \end{bmatrix} ; \text{HONEY} = \begin{bmatrix} 8 \\ 15 \\ 14 \\ 5 \\ 25 \end{bmatrix} ;
 \end{aligned}$$

Let the key matrix  $T = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$  and  $T^{-1} = \begin{bmatrix} 1 & -1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$

$$T^{-1} = \begin{bmatrix} 1 & -1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Now we multiplied the column vector corresponding to key matrix,

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 13 \\ 5 \\ 5 \\ 20 \\ 0 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 16 \\ 3 \\ 25 \\ 20 \\ 0 \end{bmatrix} = \begin{bmatrix} P \\ C \\ Y \\ T \end{bmatrix} = PCYT$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 13 \\ 5 \\ 0 \\ 1 \\ 20 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 12 \\ 26 \\ 21 \\ 21 \\ 20 \end{bmatrix} = \begin{bmatrix} L \\ Z \\ U \\ U \\ T \end{bmatrix} = LZUUT$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 0 \\ 2 \\ 21 \\ 20 \\ 20 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 9 \\ 9 \\ 7 \\ 13 \\ 20 \end{bmatrix} = \begin{bmatrix} I \\ I \\ G \\ M \\ T \end{bmatrix} = IIGMT$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 5 \\ 18 \\ 6 \\ 12 \\ 25 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 12 \\ 7 \\ 16 \\ 10 \\ 25 \end{bmatrix} = \begin{bmatrix} L \\ G \\ P \\ J \\ Y \end{bmatrix} = LGPJY$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 0 \\ 7 \\ 1 \\ 12 \\ 1 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 21 \\ 21 \\ 14 \\ 13 \\ 1 \end{bmatrix} = \begin{bmatrix} U \\ U \\ N \\ M \\ A \end{bmatrix} = UUNMA$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 24 \\ 25 \\ 0 \\ 20 \\ 15 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 3 \\ 6 \\ 8 \\ 8 \\ 15 \end{bmatrix} = \begin{bmatrix} C \\ F \\ H \\ H \\ O \end{bmatrix} = CFHHO$$



$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 0 \\ 7 \\ 5 \\ 20 \\ 0 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 5 \\ 5 \\ 25 \\ 20 \\ 0 \end{bmatrix} = \begin{bmatrix} E \\ E \\ Y \\ T \\ \end{bmatrix} = EEYT$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 8 \\ 15 \\ 14 \\ 5 \\ 25 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 13 \\ 5 \\ 17 \\ 3 \\ 25 \end{bmatrix} = \begin{bmatrix} M \\ E \\ Q \\ C \\ Y \end{bmatrix} = MEQCY$$

Hence the message to be sent is,

***PCYT LZUUTIIGMTLGPJYUUNMACFHHEEYT MEQCY***

By multiplying the inverse of key matrix T, receiver can decrypt the message easily.

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} P \\ C \\ Y \\ T \\ \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \\ 25 \\ 20 \\ 0 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 13 \\ 5 \\ 5 \\ 20 \\ 0 \end{bmatrix} \Rightarrow MEET$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} L \\ Z \\ U \\ U \\ T \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 12 \\ 26 \\ 21 \\ 21 \\ 20 \end{bmatrix} \text{mod}(27) = \begin{bmatrix} 13 \\ 5 \\ 0 \\ 1 \\ 20 \end{bmatrix} \Rightarrow MEAT$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} I \\ I \\ G \\ M \\ T \end{bmatrix} \pmod{27} = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 9 \\ 9 \\ 7 \\ 13 \\ 20 \end{bmatrix} \pmod{27} = \begin{bmatrix} 0 \\ 2 \\ 21 \\ 20 \\ 20 \end{bmatrix} \Rightarrow \text{BUTT}$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} L \\ G \\ P \\ J \\ Y \end{bmatrix} \pmod{27} = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 12 \\ 7 \\ 16 \\ 10 \\ 25 \end{bmatrix} \pmod{27} = \begin{bmatrix} 5 \\ 18 \\ 6 \\ 12 \\ 25 \end{bmatrix} \Rightarrow \text{ERFLY}$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} U \\ U \\ N \\ M \\ A \end{bmatrix} \pmod{27} = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 21 \\ 21 \\ 14 \\ 13 \\ 1 \end{bmatrix} \pmod{27} = \begin{bmatrix} 0 \\ 7 \\ 1 \\ 12 \\ 1 \end{bmatrix} \Rightarrow \text{GALA}$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} C \\ F \\ H \\ H \\ O \end{bmatrix} \pmod{27} = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 3 \\ 6 \\ 8 \\ 8 \\ 15 \end{bmatrix} \pmod{27} = \begin{bmatrix} 24 \\ 25 \\ 0 \\ 20 \\ 15 \end{bmatrix} \Rightarrow \text{XY TO}$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} E \\ E \\ Y \\ T \end{bmatrix} \pmod{27} = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 5 \\ 5 \\ 25 \\ 20 \\ 0 \end{bmatrix} \pmod{27} = \begin{bmatrix} 0 \\ 7 \\ 5 \\ 20 \\ 0 \end{bmatrix} \Rightarrow GET$$

$$\begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} M \\ E \\ Q \\ C \\ Y \end{bmatrix} \pmod{27} = \begin{bmatrix} 1 & 26 & 0 & 0 & 0 \\ 0 & 1 & 26 & 0 & 0 \\ 0 & 0 & 1 & 26 & 0 \\ 0 & 0 & 0 & 1 & 26 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 13 \\ 5 \\ 17 \\ 3 \\ 25 \end{bmatrix} \pmod{27} = \begin{bmatrix} 8 \\ 15 \\ 14 \\ 5 \\ 25 \end{bmatrix} \Rightarrow HONEY$$

Finally, we decrypted the original message **MEET ME AT BUTTERFLY GALAXY TO GET HONEY.**

### 3. Conclusion

This paper introduces the method for sending the secret messages. The key matrix and congruence modulo should be understood to decrypt the message more securely between the receiver and the sender.

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